

Oracles

Max SAT = $\{F, k \text{ st } \exists \alpha \text{ satisfying } k \text{ clauses of } F \text{ and } \forall \beta \text{ satisfies } \leq k \text{ clauses of } F\}$.

prop Max SAT $\in P^{\text{SAT}}$.

$F = C_1 \wedge C_2 \wedge \dots \wedge C_m$ original formula.

Let $G_\ell = (C_1 r r_1) \wedge (C_2 r r_2) \wedge \dots \wedge (C_m r r_m)$
 $\wedge [r_1 + r_2 + \dots + r_m \leq m - \ell]$

\hookrightarrow encode using adder circuit plus comparison circuit.

Claim: $G_\ell \in \text{SAT}$ iff F has assignment satisfying $\geq \ell$ clauses.

proof: α sat ℓ clauses in $F \Rightarrow \beta = \alpha$ extended to $r_i = [\alpha \text{ satisfies } C_i]$ has $\leq m - \ell$ r -variables set to 1 $\Rightarrow G_\ell$ satisfied.

β satisfies $G_\ell \Rightarrow \beta$ has $\leq m - \ell$ r -variables set to 1 $\Rightarrow \alpha = \beta$ restricted satisfies $\geq \ell$ clauses in F .

Hence the following oracle machine solves Max SAT in poly time:

M^{SAT} : input (F, k) .

if $G_k \in \text{SAT}$ and $G_{k+1} \notin \text{SAT}$:
 accept
 o/w:
 reject

Cook-Levin theorem.

th SAT is NP-complete.

proof. Let $L \in \text{NP}$, i.e. $L = \{x : \exists y, V(x, y) \text{ accepts}\}$.

Want to show $L \leq_p \text{SAT}$. We build a formula that simulates the computation of V .

* note it is the same formula we used in Gödel's incompleteness thm.

We'll encode that s_0, s_1, \dots, s_m is a valid computation. Have variables z_{ijk} representing state i /symbol j /bit k .

Encode: $s_0 = q_0 x y$.

$$s_{i,j-1} \quad s_{i,j} \quad s_{i,j+1}$$

$$s_{i+1,j-1} \quad s_{i+1,j} \quad s_{i+1,j+1}$$

\rightarrow If $s_{i,j-1} = q_a$ and $\delta(q_a, s_{i,j}) = (q_b, c, R)$

then $s_{i+1,j-1} = c$ and $s_{i+1,j} = q_b$

\rightarrow If $s_{i,j} = q_a$ and $\delta(q_a, s_{i,j+1}) = (q_b, c, L)$

then $s_{i+1,j+1} = c$ and $s_{i+1,j} = q_b$

\rightarrow O/w $s_{i+1,j} = s_{i,j}$

Obs we are encoding a condition over constant many variables. Hence exists CNF of constant size.

We need $|s_i| = \text{poly}(n)$ checks per state and we have $m = \text{poly}(n)$ many states. Hence total size of CNF is $\text{poly}(n)$.

This proves reduction is poly size. Correctness is easy to check from definitions.