

Cumulative Space in Black-White Pebbling and Resolution

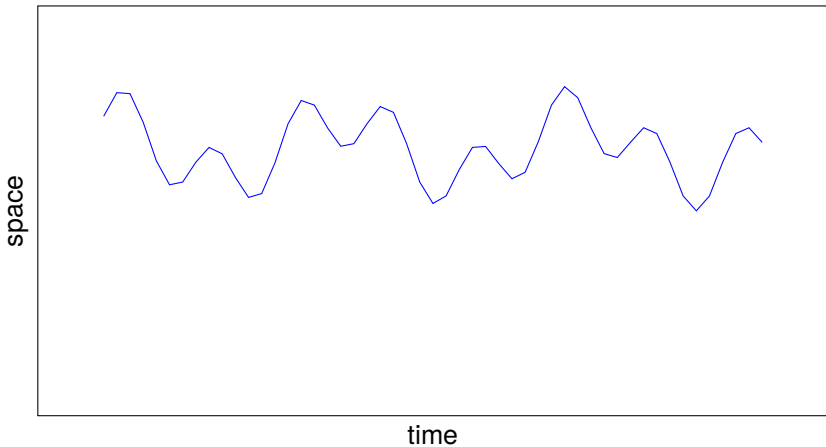
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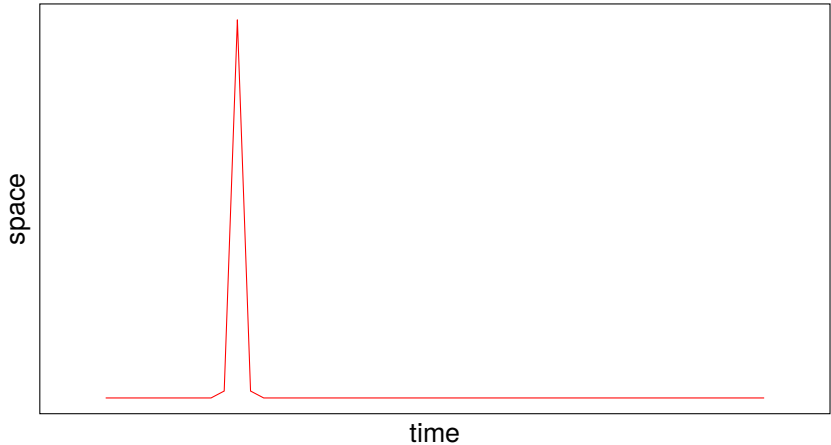
joint work with Joël Alwen (IST Austria),
Susanna F. de Rezende (KTH),
and Jakob Nordström (KTH)

8th Innovations in Theoretical Computer Science

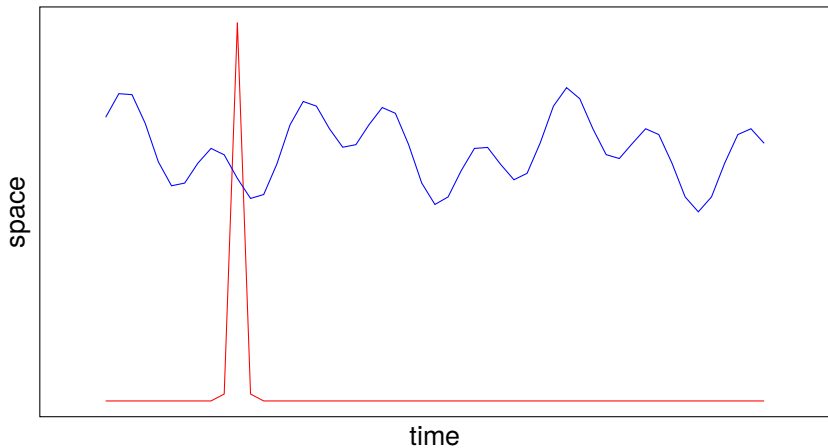
What is space?



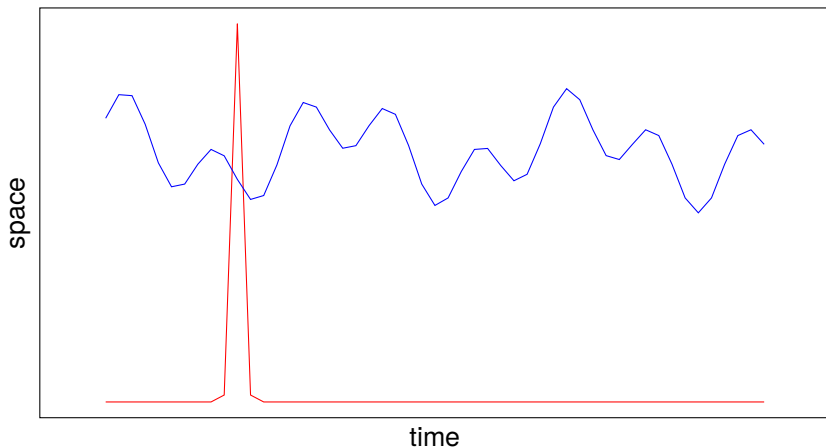
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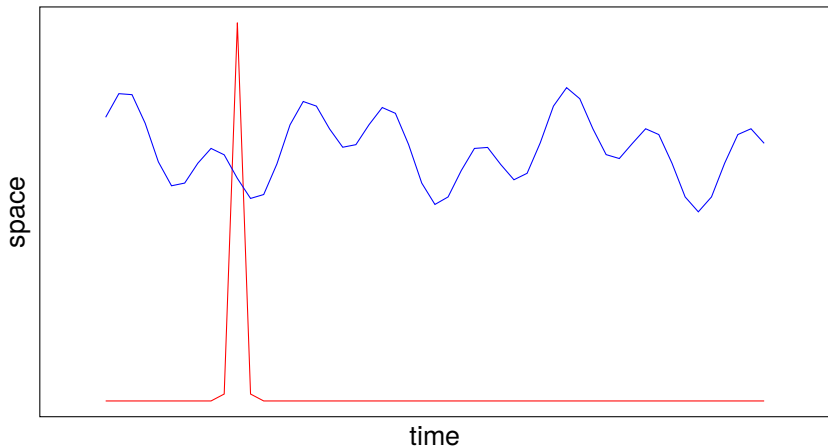


What is space?



Usually: *maximal* space.

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[Alwen, Serbinenko '15]: aggregate space over computation (*cumulative* space).

Resolution

Setup

Prove CNF formula unsatisfiable.

Present proof on board.

▶ Write down axiom clauses

▶ Infer new clauses

$$\frac{C \vee x \quad D \vee \bar{x}}{C \vee D}$$

▶ Erase clauses to save space

Goal: derive empty clause \perp

$$F = \{x, \bar{x} \vee y, \bar{y}\}$$



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$$\frac{\bar{x} \vee y}{y}$$

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Questions

▶ How much time will this take? (Length)

▶ How large is the blackboard? (Space)

$$F = \{x, \bar{x} \vee y, \bar{y}\}$$

$$\begin{array}{c} y \\ \bar{y} \\ \perp \end{array}$$

Space

[Esteban, Torán '99]

[Alekhnovich, Ben Sasson, Razborov, Wigderson '00]

x	x $\bar{x} \vee y$	x $\bar{x} \vee y$ y	x $\bar{x} \vee y$ y	$\bar{x} \vee y$ y	y \bar{y}	y \bar{y} \perp
-----	-------------------------	--------------------------------	--	---	------------------	-----------------------------

$$|\mathbf{C}_1| = 1 \quad |\mathbf{C}_2| = 2 \quad |\mathbf{C}_3| = 3 \quad |\mathbf{C}_4| = 2 \quad |\mathbf{C}_5| = 1 \quad |\mathbf{C}_6| = 2 \quad |\mathbf{C}_7| = 3$$

Space of a proof: $\text{Sp}(\pi) := \max_t |\text{Clauses in } \mathbf{C}_t| = 3$

Space of refuting a formula: $\text{Sp}(F \vdash \perp) := \min_{\pi: F \vdash \perp} \text{Sp}(\pi) \leq 3$

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Alternative measures: # literals, # bits

Space

Bounds

Every formula $Sp = O(n)$

Exist formulas st $Sp = \Omega(n)$

[Esteban, Torán '99], [Alekhnovich, Ben Sasson, Razborov, Wigderson '00]

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Space vs length

Exist formulas st

- ▶ Exists proof with $Sp = O(n^{1/11})$
- ▶ Exists proof with $Len = O(n)$
- ▶ Every proof with $Sp < n^{2/11}$ requires $Len = \exp n^{\Omega(1)}$

[Ben Sasson, Nordström '11]

Cumulative Space

Aggregate space over whole proof.

x	x $\bar{x} \vee y$	x $\bar{x} \vee y$ y	x $\bar{x} \vee y$ y	$\bar{x} \vee y$ y	y \bar{y}	y \bar{y} \perp
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Cumulative space of a proof: $\text{CumSp}(\pi) := \sum_t |\text{Clauses in } \mathbf{C}_t| = 14$

Cumulative space of refuting a formula:

$$\text{CumSp}(F \vdash \perp) := \min_{\pi: F \vdash \perp} \text{CumSp}(\pi) \leq 14$$

Cumulative Space

Observations

Every proof $\text{CumSp} \leq \text{Len} \cdot \text{Sp}$

Every formula $\text{Len} \leq 2^n$ and $\text{Sp} \leq n$

$$\Rightarrow \text{CumSp} \leq n2^n$$

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Every formula $\text{CumSp} \leq \text{Len}^2$.

Reaching space s needs $s/2$ configurations of space $\geq s/2$

\Rightarrow Cumulative space $\Omega(s^2)$.

Cumulative Space Bounds

How large can cumulative space be?

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Maximal space: $\text{Sp} = O(\text{Len})$ not tight.

Every formula $\text{Sp} = O(\text{Len} / \log \text{Len})$. [Hopcroft, Paul, Valiant '75]

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Theorem

Exist formulas with $\text{Len} = O(n)$ and $\text{CumSp} = \Omega(n^2)$.

Maximal vs Cumulative Space

Large space \Leftrightarrow large cumulative space?

\Rightarrow Yes

Every formula $\text{CumSp} = \Omega(\text{Sp}^2)$.

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\Leftarrow No

Theorem

Exist formulas with $\text{Sp} = O(\log n)$ but $\text{CumSp} = \Omega(n^2 / \log n)$

Length vs Cumulative Space

How often do we need maximum space in a trade-off?

Theorem [Ben Sasson, Nordström '11]

Exist formulas st for any $s = O(\sqrt{n})$

- ▶ Exists proof with $Sp = O(s)$ and $Len = O(n^2/s^2)$
 - ▶ Exists proof with $Sp = O(1)$
 - ▶ Exists proof with $Len = O(n)$
- ▶ Every proof in space $O(s)$ needs $Len = \Omega(n^2/s^2)$

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 - ▶ Exists proof with $Sp = O(1)$
 - ▶ Exists proof with $Len = O(n)$
- ▶ Every proof in space $O(s)$ needs $CumSp = \Omega(n^2/s)$

Corollary

- ▶ Every proof in space $O(s)$ and length $O(n^2/s^2)$ needs $\Omega(n^2/s^2)$ configurations with space $\Omega(s)$

Parallel Resolution

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Automatic $\text{CumSp} = \Omega(\text{Sp}^2)$ lower bound no longer holds.

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Previous results hold even allowing parallel inference.

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Fully Parallel Resolution

Very powerful model: can prove any formula in 2 steps.

Lower bounds with limited space.

Techniques

Pebble games

- ▶ Simple computational model to measure space.
- ▶ Prove lower bounds in pebble game
- ▶ Translate to resolution

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Lemma

Resolution proof of $F(G)$ in length L , space s , cumulative space c .
Then pebbling of G in time L , space s , cumulative space c .

Even if parallel inference steps.

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Even if parallel inference steps.

- ▶ [Alwen, Serbinenko '15]: Translate computation to black pebbling strategy.
- ▶ Proofs are non-deterministic: translate proof to black-white pebbling.

Take Home

Recap

- ▶ Introduced cumulative space measure in proof complexity.

Open problems

- ▶ Study cumulative space in other areas.

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Thanks!